(54) Title: TIME SYNCHRONIZATION IN COMPUTER NETWORK

(57) Abstract: A method for improved Network Time Protocol time synchronization in a computer network (1) where the computer
network comprises a time client (2) and a timeserver (3). To eliminate network access errors at the client side, the actual time of
leaving is stored in the time client (2) as a special time stamp (T11) and later substituted for T1 in the calculations. To eliminate
network access errors at the server side, the time request packet is duplicated and returned twice, the second time containing the time
stamp when the first packet left the server (T31). To eliminate network transversal jitter when using a switched network, a packet
with a multicast is used and the time stamp (T11) is taken when the time request packet is reflected back from the switch and (T31)
is taken when the reply packet is reflected back from the switch.

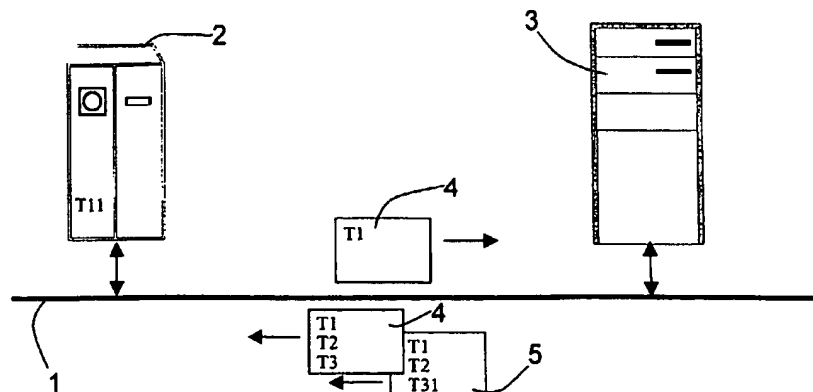MK, MN, MW, MX, MZ, NO, NZ, OM, PH, PL, PT, RO, RU, SD, SE, SG, SI, SK, SL, TJ, TM, TN, TR, TT, TZ, UA, UG, UZ, VN, YU, ZA, ZM, ZW, ARIPO patent (GH, GM, KE, LS, MW, MZ, SD, SL, SZ, TZ, UG, ZM, ZW), Eurasian patent (AM, AZ, BY, KG, KZ, MD, RU, TJ, TM), European patent (AT, BE, CH, CY, DE, DK, ES, FI, FR, GB, GR, IE, IT, LU, MC, NL, PT, SE, TR), OAPI patent (BF, BJ, CF, CG, CI, CM, GA, GN, GQ, GW, ML, MR, NE, SN, TD, TG)

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For two-letter codes and other abbreviations, refer to the "Guidance Notes on Codes and Abbreviations" appearing at the beginning of each regular issue of the PCT Gazette.

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TIME SYNCHRONIZATION IN COMPUTER NETWORK

Field of the invention

5 A method for enhanced accuracy Network Time Protocol time synchronization in a computer network, such as a local area network comprising a time client and a timeserver.

Background of the invention

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The IETF (Internet Engineering Task Force group) Network Time Protocol (NTP) standard RFC 1305 defines a method for synchronizing workstation clocks across the Internet. This method has an accuracy of about 1 ms (millisecond), which is adequate for time stamping files and other non-real-time operating system chores. Certain classes of automation systems, the most notable being Substation Automation, i.e. the control and protection of energy distribution network nodes, require much more precise time synchronization, for example, 1µs for class 1 applications and 25µs for class 2 applications. A class 1 application is, for example, time tagging of syncrophasors and a class 2 application is, for example, time tagging of phasors.

25 The RFC 1305 standard comprises an algorithm for calculation of the corrections to a time-of-day clock in one node, such as a time client, relative to a reference time-of-day clock in another node, such as a timeserver. The algorithm is based on a network packet, from now on

30 called a time-request packet, containing three important time stamps:

T1 (*Originate Timestamp*): The time the time request packet was generated in the client asking for the current time.

35 T2 (*Receive Timestamp*): The time the time request packet arrived at the timeserver.

T3 (*Transmit Timestamp*): The time the time request packet was updated and put into a transmission queue at the timeserver.

40 In addition, the calculations require:

T4: The time the time request packet arrived back at the time client.

T2 and T4 are easily determined with accuracy down to
5   microseconds, and sometimes even more accurate, using hardware or software time stamps based on network packet arrival interrupts.
Accurate determination of T1 and T3 is, however, a problem. For full accuracy, T1 and T3 should be the time
10   when the network packet leaves the time client or the timeserver. The problem is that T1 and T3 are not available until the network packet has already left the timeserver or time client and then it is too late to incorporate them into the packet. Therefore, the largest
15   part of the time synchronization inaccuracy for an NTP setup is the variation in the delay between T1 and the actual time the network packet leaves the time client, as well as the variation in the delay between T3 and the actual time the network packet leaves the time server.
20

Summary of the invention

The object of the invention is to provide a method for Network Time Protocol (NTP) or Simple Network Time
25   Protocol (SNTP) time synchronization in a computer network, without the disadvantages mentioned under background of the invention.

This object is achieved by a method according to the
30   independent claim 1.

The invention provides a method where the accuracy of the time stamps involved may be substantially increased while still being compatible with the original protocol.
35
In one embodiment of the invention, the computer network is a Local Area Network (LAN) and uses the Internet communications protocol suite, usually denoted TCP/IP.

40   According to another preferred embodiment, the invention

is used in control and protection of an energy
distribution network node for improving time accuracy in
the Network Time Protocol.

5   According to another preferred embodiment of the
invention, the Network Time Protocol is extended in a
backwards-compatible way such that accuracy in the order
of 10µs or better may be attained.

10 Brief description of the drawing

Fig. 1 is a schematic block diagram of a computer
network, such as a LAN, comprising a time client and a
timeserver where a time request packet is transmitted
15 between the time client and the timeserver.

Fig. 2 is a schematic block diagram of a computer
network, such as a LAN, with a store-and-forward and/or
switching device arranged between the client and the
20 server.

Detailed description of preferred embodiments

Figure 1 shows a schematic block-diagram of a computer
25 network 1, such as a Local Area Network (LAN), comprising
a time client 2 and a timeserver 3. A time request packet
4 is transmitted between the time client and the
timeserver.

30 The computer network protocol used is the Internet
communication suite, usually denoted TCP/IP.

The two main sources of inaccuracy that plagues the
standard NTP method are reduced with the method described
35 below. The method is concerned with a first and a second
part of the algorithm in the standard NTP method, which
will be described in the following.

The algorithm is based on the time-request packet,
40 containing three important time stamps:

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A first time stamp T1: The time the time request packet
was generated in the client asking for the current time.
A second time stamp T2: The time the time request packet
arrived at the timeserver.

5   A third time stamp T3: The time the time request packet
was updated and put into a transmission queue at the
timeserver.

The first main source of inaccuracy is the variation in
10   the delay between the first time stamp T1 and the exact
time when the time request packet leaves the time client.
To improve the time client accuracy the following steps
are performed:

1. Create an NTP time request packet 4 and fill in the
15   first time stamp T1.
2. Transmit the time request packet to the timeserver
   3.
3. Get hold of the actual time stamp, hereinafter
   called the fourth time stamp T11, when the time
20   request packet leaves the time client 2, using the
   network transmit interrupt or a hardware time stamp.
4. Store the fourth time stamp T11 in a data structure
   in the time client together with the first time
   stamp T1.The timeserver 3 receives the time request
25   packet 4 and transmits it back to the time client 2
5. Replace the first time stamp T1 in the time request
   packet by the fourth time stamp T11, when the time
   request packet 4 returns from the timeserver.,

This exchange eliminates the inaccuracy associated with
30   the originate timestamp T1 in the time calculation and
thus an improved NTP time client is achieved. This
corresponds to the first part in the algorithm.

The second main source of inaccuracy is the variation in
35   the delay between the third time stamp T3 and the actual
time the packet leaves the timeserver 3. To improve the
timeserver accuracy, the following steps are preformed:

1. Create a duplicate packet 5 before entering the
   third time stamp T3 into the time request packet,
40   when the time request packet arrives at the

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timeserver 3. This packet will hereinafter be
denoted the time correction packet 5.
2. Transmit the time request packet 4 back to the time
client 2.
3. Acquire an accurate time stamp, hereinafter called
the fifth time stamp T31, when the time request
packet leaves the timeserver, using the network
transmission interrupt or a hardware time stamp.
4. Put the fifth time stamp T31 into the third time
stamp T3 location in the time correction packet.
5. Transmit the time correction packet back to the time
client.

Now, the returned time request packet 4 has the standard
NTP inaccuracy in the third time stamp T3. However, if
the third time stamp T3 in the returned time request
packet 4 is replaced by the fifth time stamp T31 from the
time correction packet 5, the accuracy of the time stamps
in the resulting packet will be:
   • T1: Maximum if the substitution described in the
     first part of the algorithm is performed.
   • T2: Maximum always.
   • T3: Maximum if the substitution described in the
     second part of the algorithm is performed.
In this way network access jitter is eliminated and only
network transmission jitter is left.

Compatibility with Network Time Protocol
If an improved NTP time client connects to a standard NTP
timeserver, only one packet will be returned, i.e. the
time request packet. The improved time client executes
the first part of the algorithm eliminating the first
time stamp T1 inaccuracy. The resulting accuracy will
therefore be better than standard NTP.

If a standard NTP client connects to an improved NTP time
server, it will get two time packets from the time
server, i.e. the time request packet and the time
correction packet, both of these having standard NTP
accuracy. If the first packet is received successfully,

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the time will be updated according to the contents of
that packet and the second packet will be discarded as an
unnecessary duplicate. If the first packet is lost, the
time will be updated according to the contents of the
5   second packet. The second packet contains the T31
timestamp which, taken by itself, is a T3 timestamp with
standard NTP accuracy.

Reducing network transmission jitter
10   In the case of a standard Local Area Network, the
difference between the time stamp taken when the time
request packet 4 leaves the client (server) transmitter
and the time stamp taken when the same packet arrives at
the server (client) receiver varies very little between
15   transmissions. In the case where a store-and-forward
and/or switching device 6 are/is in the path between the
client and the server as in switched Ethernet, this
difference will vary according to the traffic load at the
time. Therefore, a store-and-forward and/or switching
20   device 6 will introduce an unpredictable jitter in the
packet travel time between transmitter and receiver.

One way of reducing this jitter is to use broadcast or
multicast packets as time request packets, between the
25   time client 2 and the timeserver 3, together with a full
duplex link to the store-and-forward device and/or
switching device. The store-and-forward device and/or
switching device will then be forced to forward this
multicast packet to all connected nodes, usually with a
30   time difference below one microsecond. In particular, the
packet will be sent back to the originator, i.e. the time
client or server, at a time very close to the time it is
sent to the destination, i.e. the timeserver or client.
This feature enables the originator to time stamp the
35   reflected packet with a time stamp closely connected to
the time the packet was sent to the destination from the
store-and-forward device, eliminating the store-and-
forward jitter.

40   Figure 2 shows an example of how to reduce the network

transmission jitter when the store-and-forward device
and/or switching device 6 is in the path between the
timeserver 2 and the time client 3. The time request
packet 4 is generated in the time client, and is stamped
5   with a first time stamp T1 corresponding to the time it
is generated. The time request packet is transmitted to
the timeserver using a multicast or broadcast address. On
the way to the timeserver, the time request packet passes
the store-and-forward device and/or switching device, and
10   when the request packet is reflected back to the time
client from the store-and-forward and/or switching
device, it is stamped with the fourth time stamp T11.
This fourth time stamp T11 is stored in a data structure
in the client together with the first time stamp T1.
15
The time request packet that is transmitted to the
timeserver is stamped with the second time stamp T2 when
it arrives at the timeserver 3. Now, the time correction
packet 5 is created before entering the third time stamp
20   T3 into the time request packet. The time request packet
is sent back to the time client from the timeserver,
using a multicast or broadcast address. On the way back
to the time client, the time request packet passes the
store-and-forward device and/or switching device, and is
25   reflected back to the timeserver where it is stamped with
a fifth time stamp T31. This fifth time stamp is stored
in the third time stamp T3 location in the time
correction packet 5 and transmitted to the time client.

30   The time client is able to do the same substitutions as
described before, i.e. substituting T11 for T1 and T31
for T3, before calculating the time correction. The
difference is that also the store-and-forward delay
variations from the calculations are eliminated, further
35   improving the accuracy.

In a preferred embodiment of the invention the time
stamping of the fourth time stamp T11 or the fifth time
stamp T31 is performed either in a network transmit
40   interrupt or by using a dedicated hardware timer.

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Before the calculations of the correct time, the time
request packet 4 is also time-stamped with a sixth time
stamp T4 (not shown) as it arrives back at the time
client.

Claims

1.   A method for Network Time Protocol time
synchronization in a computer network (1), the computer
network comprising a time client (2) and a timeserver
(3), the method comprising the steps of
-    generating a time request packet (4) in the time
client,
-    time-stamping the time request packet with a first
time stamp (T1) corresponding to the time the time
request packet is generated,
-    transmitting the time request packet to the time
server,
-    time-stamping the time request packet with a second
time stamp (T2) when it arrives at the time server,
-    time-stamping the time request packet with a third
time stamp (T3) when it is sent back to the time client,
c h a r a c t e r i z e d   b y   the steps of
-    storing the actual time the time request packet (4)
leaves the time client (2) as a fourth time stamp (T11)
in the time client (2), after the step of time stamping
the packet with the first time stamp (T1) in the time
client, and
-    replacing the first time stamp (T1) in the time
request packet by the fourth time stamp (T11) when the
time request packet has returned to the time client, in
order to improve the time client accuracy.

2.   A method according to claim 1,
c h a r a c t e r i z e d   i n   that the computer
network (1) is a Local Area Network (LAN).

3.   A method according to claim 1 or 2,
c h a r a c t e r i z e d   b y the steps of
-    duplicating the time request packet (4) into a time
correction packet (5) in the time server (2),
-    time-stamping the time correction packet with the
actual time the time request packet leaves the time
server in the form of a fifth time stamp (T31),

- transmitting the time correction packet (5) back to the time client (2), and
- replacing the third time stamp (T3) in the time request packet by the fifth time stamp (T31) from the

5   time correction packet.

4. A method according to claim 1, where a store-and-forward device and/or a switching device (6) are/is in the path between the time client (2) and the timeserver

10   (3), c h a r a c t e r i z e d b y
- using a multicast or broadcast address for the time request packet (4) when transmitting the time request packet to the time server such that the time request packet is reflected to the time server from the store-

15   and-forward device and/or switching device, and
- time-stamping the fourth time stamp (T11) in the time client (2) as the reflected time request packet returns to the time client (2) from the store-and-forward or switching device (6).

20

5. A method according to claim 4, c h a r a c t e r i z e d b y the steps of
- duplicating the time request packet (4) into a time correction packet (5) in the time server (2),

25   - time-stamping the time request packet with a third time stamp (T3) when it is sent back out using a multicast or broadcast address
- time-stamping the time correction packet with the actual time the time request packet was reflected back to

30   the time server in the form of a fifth time stamp (T31)
- transmitting the time correction packet back to the time client (2) and
- replacing the third time stamp (T3) in the time request packet by the fifth time stamp (T31) from the

35   time correction packet.

6. A Method according to any of the preceding claims, c h a r a c t e r i z e d b y time-stamping the fourth time stamp (T11) or the fifth time stamp (T31), either in

40   a network transmit interrupt or by using a dedicated

hardware timer.

7. Use of the method according to any of claims 1-6 for
control and protection of an energy distribution network
node.

Fig 1

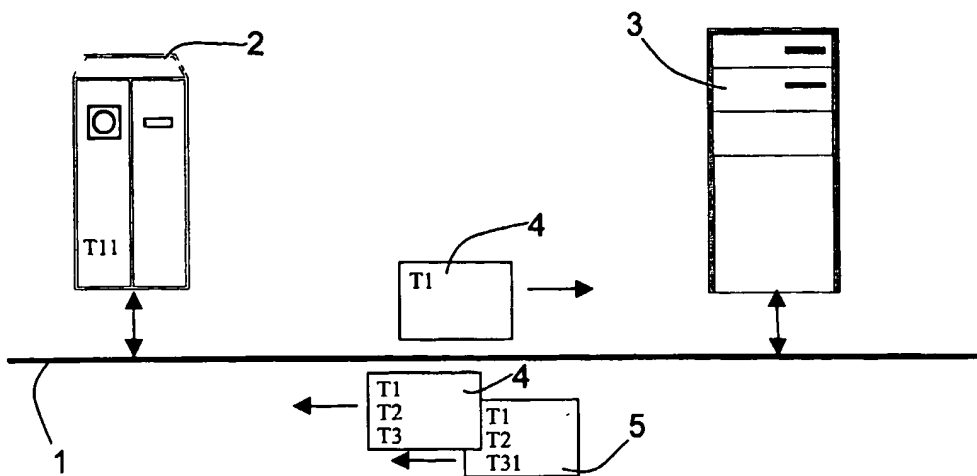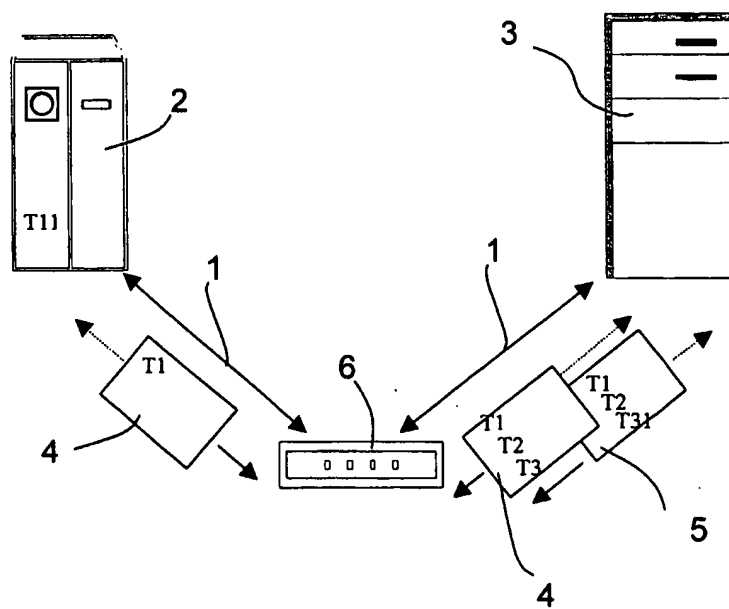
Fig 2

INTERNATIONAL SEARCH REPORT

International application No.
PCT/NO 01/00516

A. CLASSIFICATION OF SUBJECT MATTER

IPC7: H04L 7/00, H04L 12/54

According to International Patent Classification (IPC) or to both national classification and IPC

B. FIELDS SEARCHED

Minimum documentation searched (classification system followed by classification symbols)

IPC7: H04L, H04J, H04Q

Documentation searched other than minimum documentation to the extent that such documents are included in the fields searched

SE,DK,FI,NO classes as above

Electronic data base consulted during the international search (name of data base and, where practicable, search terms used)

EPO-INTERNAL, WPI-DATA, INSPEC

C. DOCUMENTS CONSIDERED TO BE RELEVANT

Category*	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
P,X	WO 0195562 A2 (ONTIME NETWORKS AS), 13 December 2001 (13.12.01), figures 1-8, claims 1-15, abstract, see whole document	1-2
A	EP 0613271 A2 (HONEYWELL INC), 31 August 1994 (31.08.94), column 1, line 1 - column 2, line 33; column 8, line 6 - line 37, figures 6,8, claims 1-5, abstract	1-7
A	TrueTime. Precise synchronization of Computer Networks: Network Time Protocol (NTP) for TCP/IP. 05-07-1997. Retrieved on 26-03-2002 from the Internet: http://www.truetime.com/DOCSn/ap23.pdf See whole document	1-17

X	Further documents are listed in the continuation of Box C.		X	See patent family annex.

* Special categories of cited documents:

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Date of the actual completion of the international search	Date of mailing of the international search report
27 March 2002	0 9 -04- 2002

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INTERNATIONAL SEARCH REPORT

	International application No.
	PCT/NO 01/00516

C (Continuation). DOCUMENTS CONSIDERED TO BE RELEVANT

Category*	Citation of document, with indication, where appropriate, of the relevant passages	Relevant to claim No.
P,A	HOLMEIDE,Ø. ET AL: Time Synchronization in Switched Ethernet. September 2001. Retrieved on 26-03-2002 from the Internet: http://ethernet.industrial-networking.com/articles/i07switched.asp See whole document -- --------	1-7

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WO	0195562	A2	13/12/01	NO	20002883	A	07/12/01
EP	0613271	A2	31/08/94	DE	69425167	D,T	22/03/01
				JP	7311721	A	28/11/95
				US	5384563	A	24/01/95